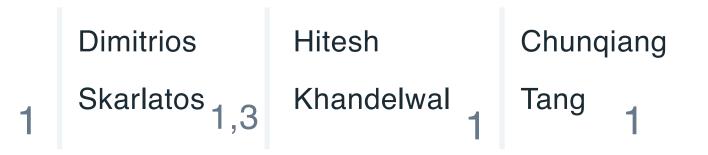
ServiceRouter

HYPERSCALE AND MINIMAL COST SERVICE MESH AT META

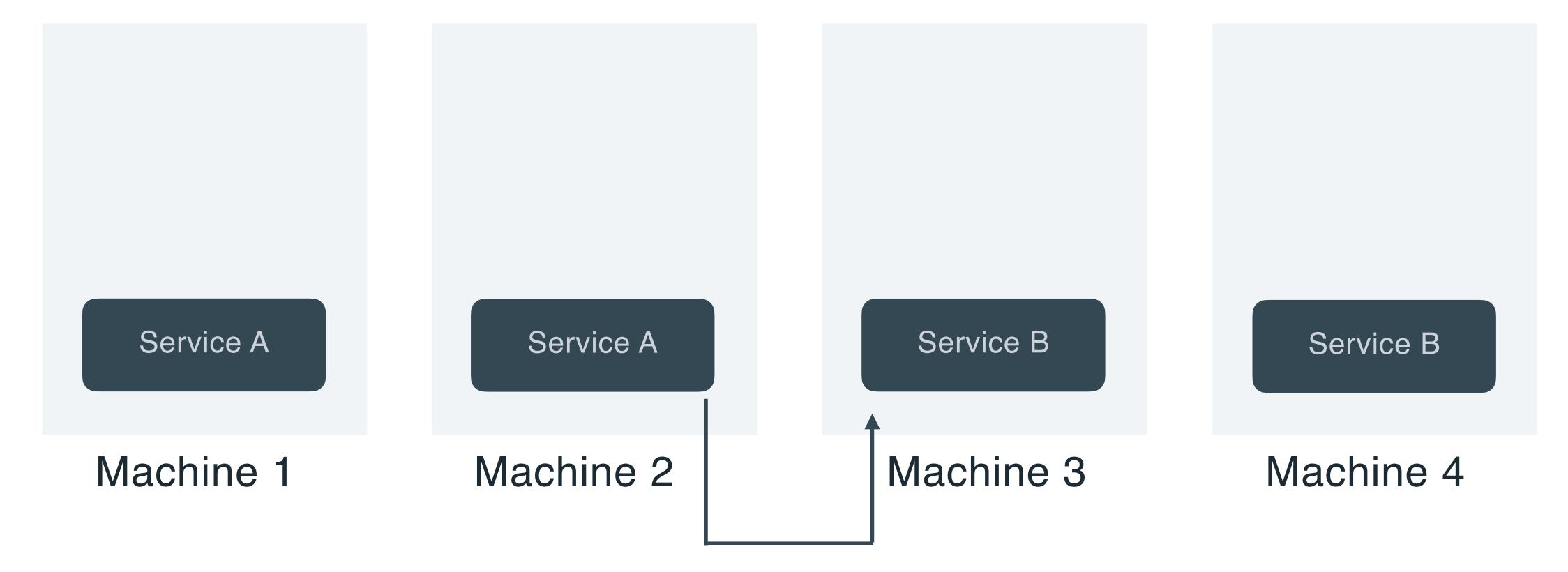
Harshit	Soteris	Nick	Max	Josh	Margot
Saokar <mark>1</mark>	Demetriou 1,2	Magerko 1	Kontorovich 1	Kirstein 1	Leibold





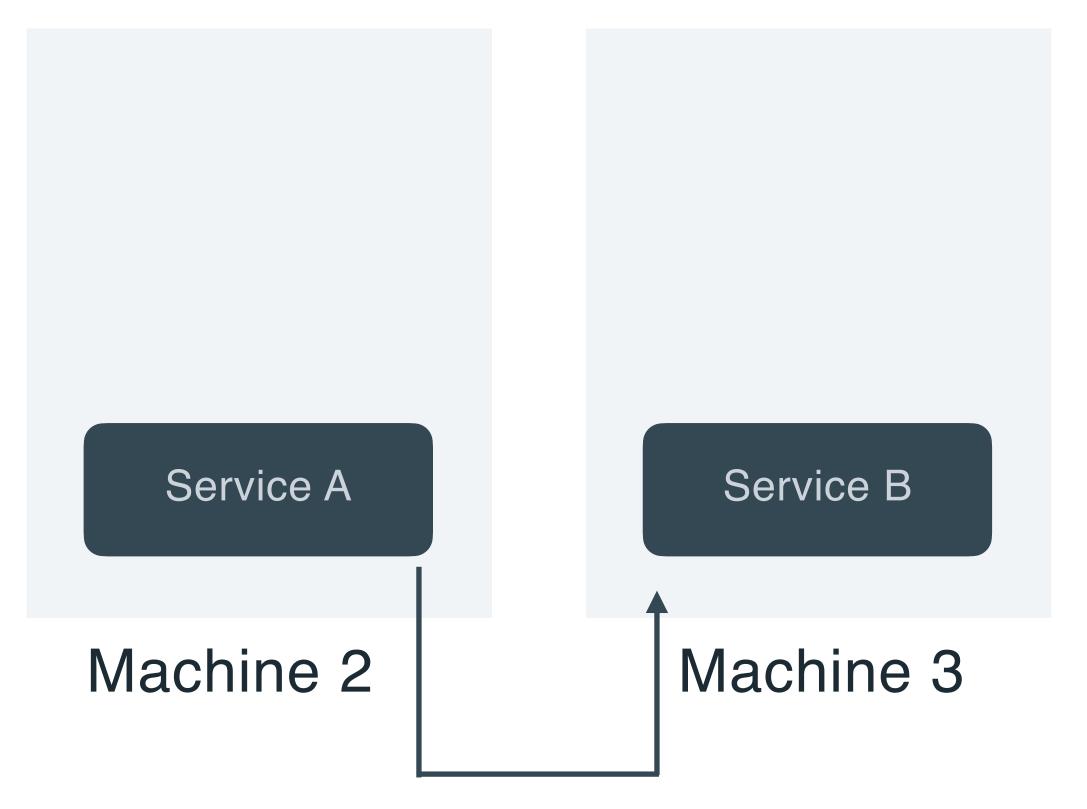
01 Background & Motivation

01 Background & Motivation



Huye et al. Lifting the veil on Meta's microservice architecture: Analyses of topology and request workflows. USENIX ATC '23

01 Background & Motivation



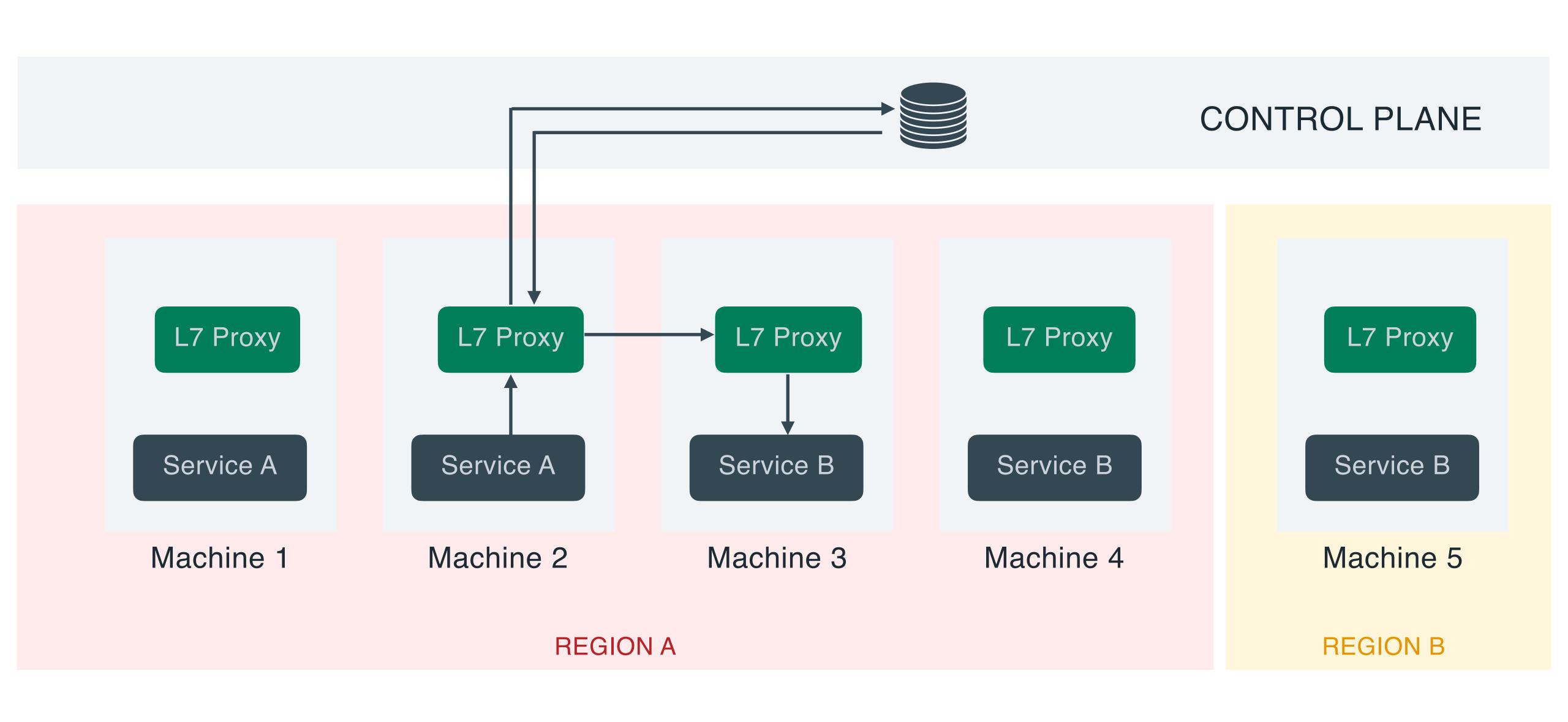
Service B

Machine 4

RPC Frameworks

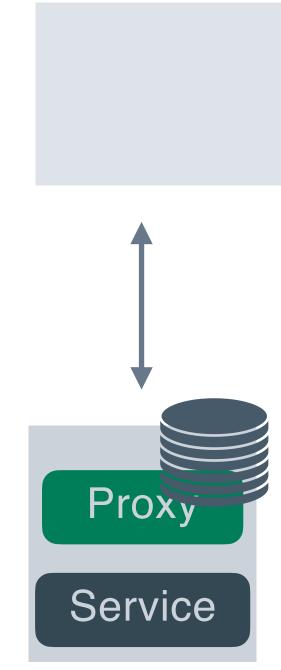
- No Advanced Load Balancing
- Need external support for service discovery
- Examples: gRPC, Thrift

Background & Motivation: Service Mesh

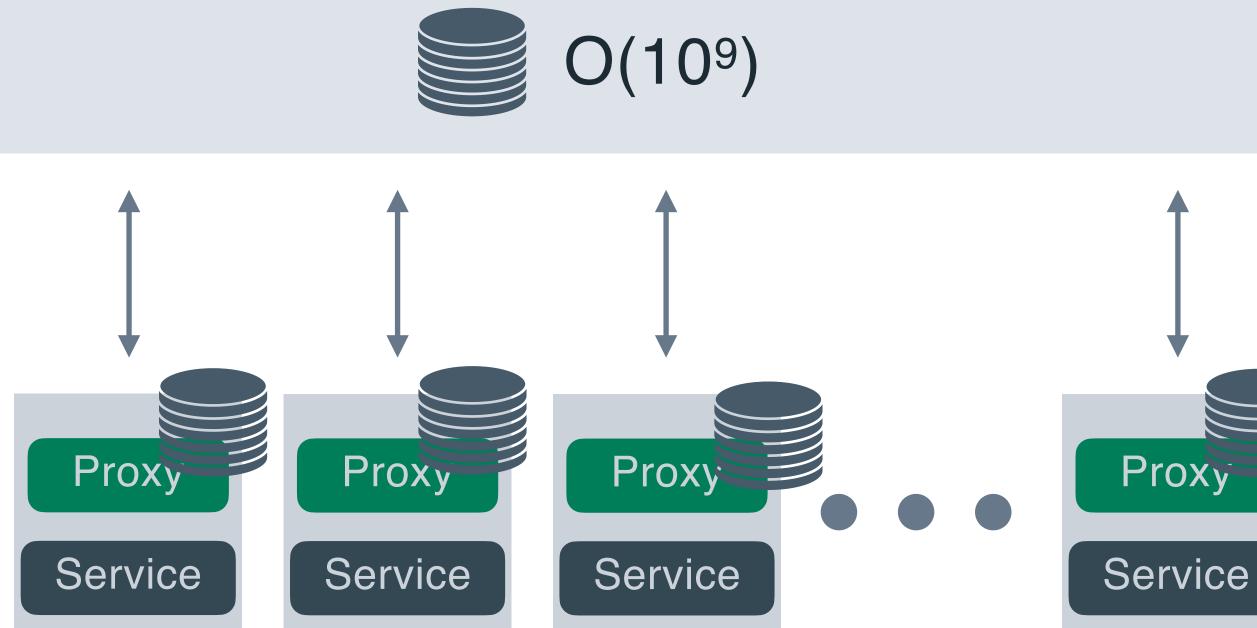


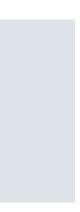
Service Mesh Challenges

• [SCALABILITY] How can we scale service discovery to O(10⁶) clients and proxies?



$O(10^{6})$





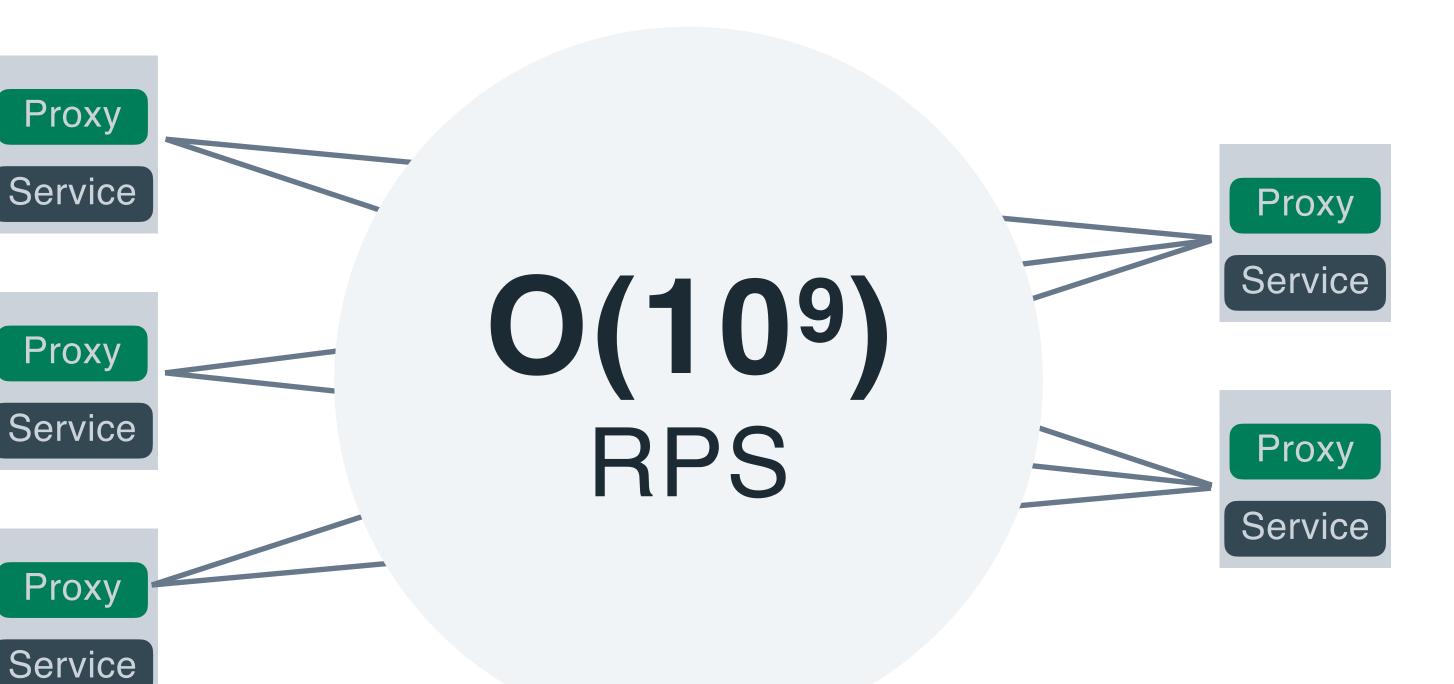


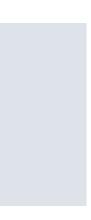
Service Mesh Challenges

- [SCALABILITY] How can we scale service discovery to O(10⁶) clients and proxies?
- **[HW COST]** How to minimize HW cost?

Istio: 0.35vCPU for O(10³) rps

1,750,000 AWS t4g.small VMs for 10B rps





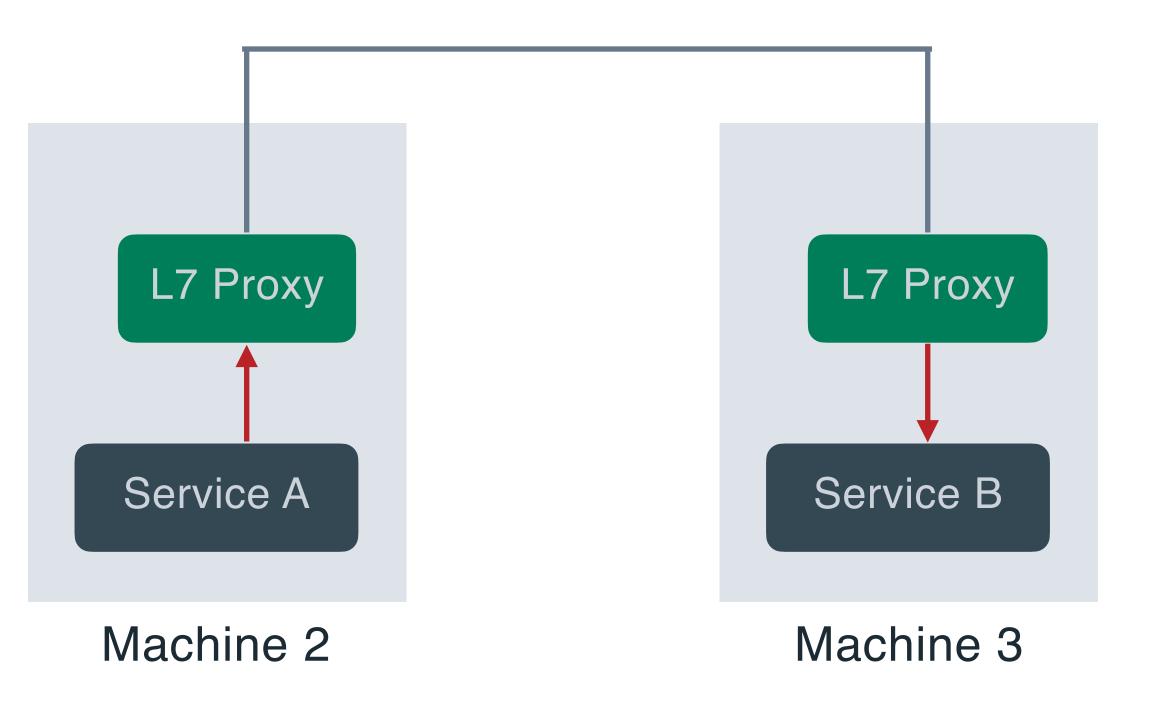
Service Mesh Challenges

- [SCALABILITY] How can we scale service discovery to O(10⁶) clients and proxies?
- [HW COST] How to minimize HW cost?
- [RPC LATENCY & LB] How to • simultaneously minimize RPC latency and load balance across geo-distributed hosts?
 - Sidecars add extra latency

Zhu et al show that Istio



Zhu et al. Dissecting Service Mesh Overheads. In *arXiv preprint arXiv:2207.00592*, 2022.



increases the latency by 185%

mRPC shows that a sidecar approach:

increases P99 RPC latency by 180%



Chen, et al. Remote procedure call as a managed system service. NSDI '23

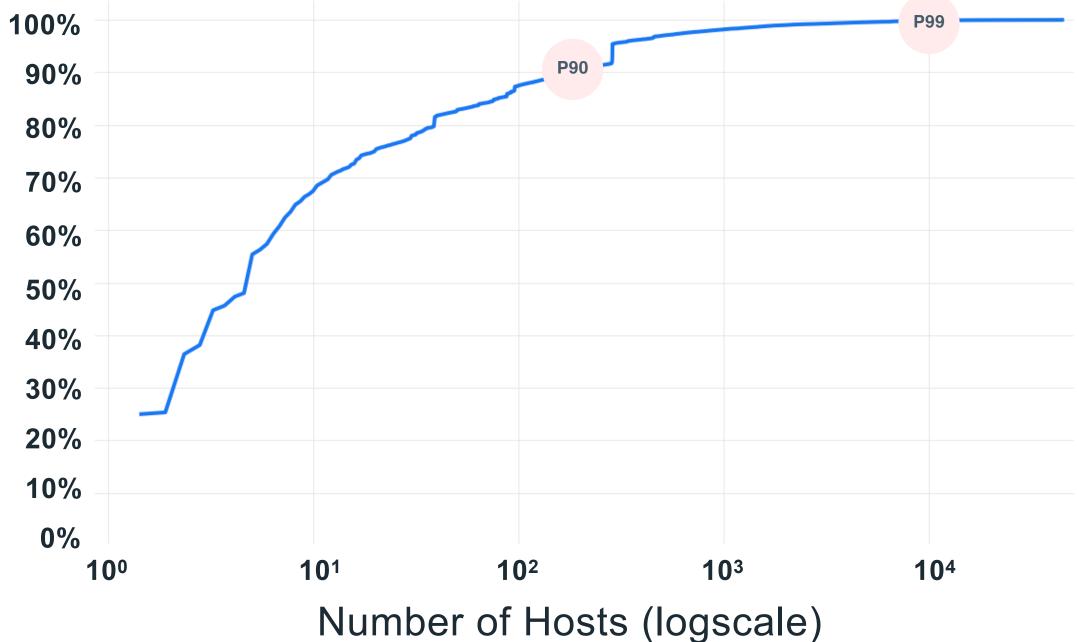


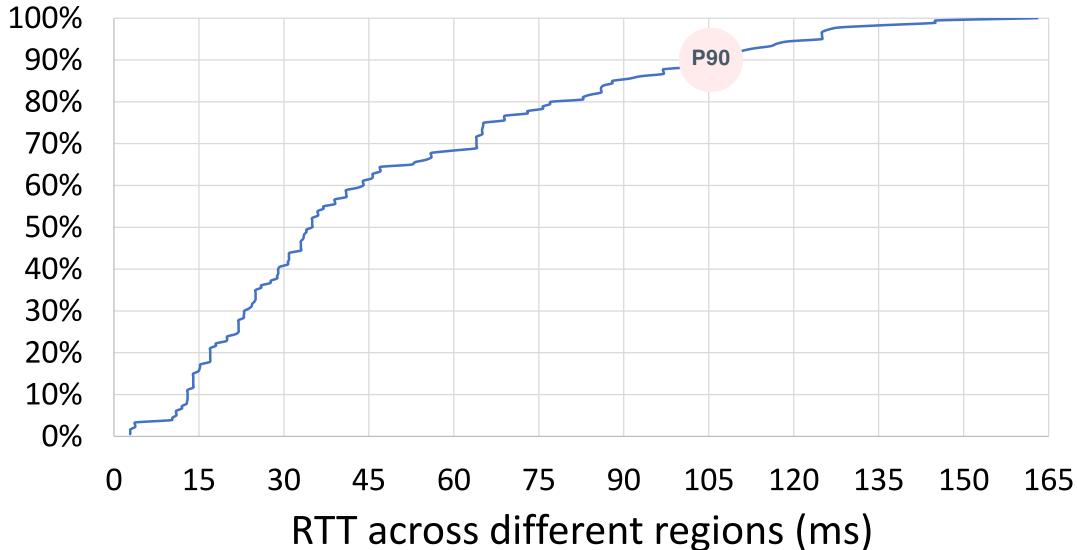
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 - Sidecars add extra latency
 - O(10-10⁴) hosts per service
 - P90 cross-region latency: 106ms

CDF

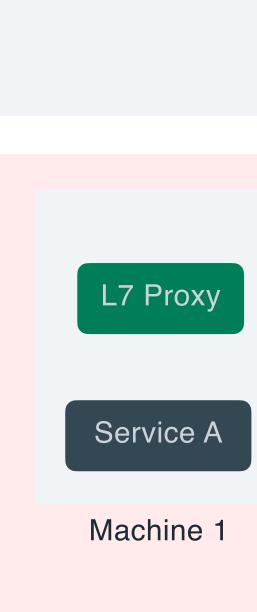
CDF

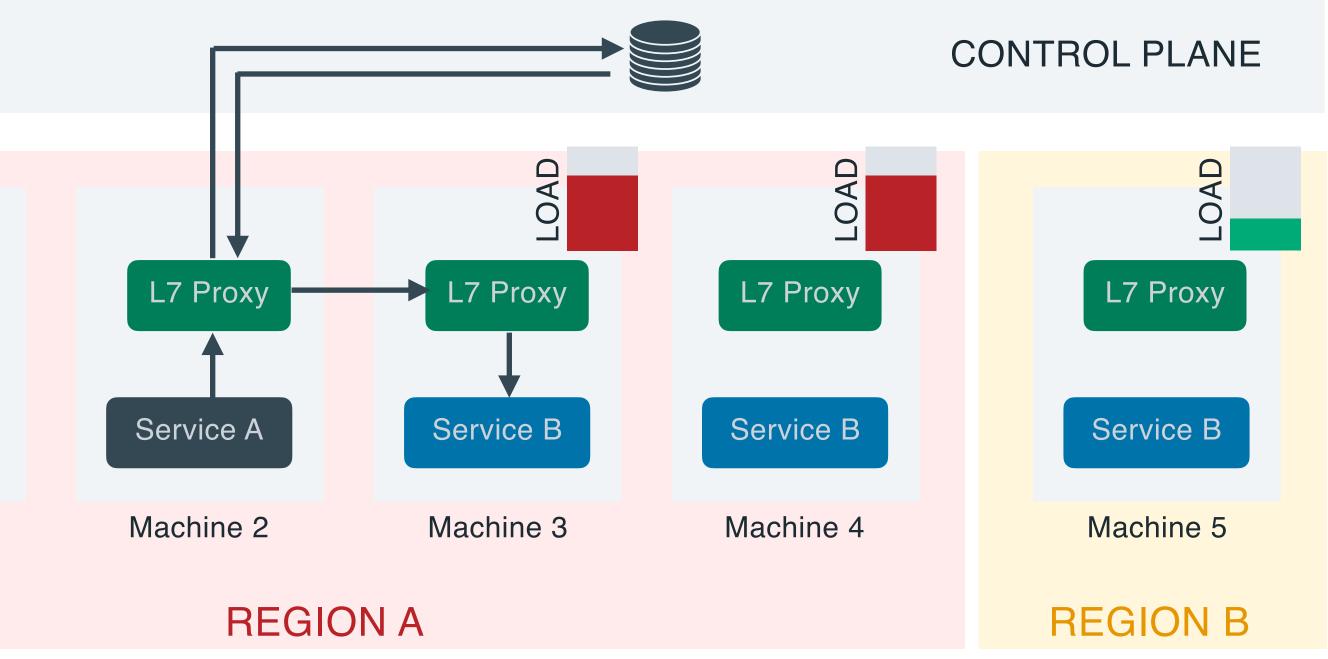




Service Mesh Challenges

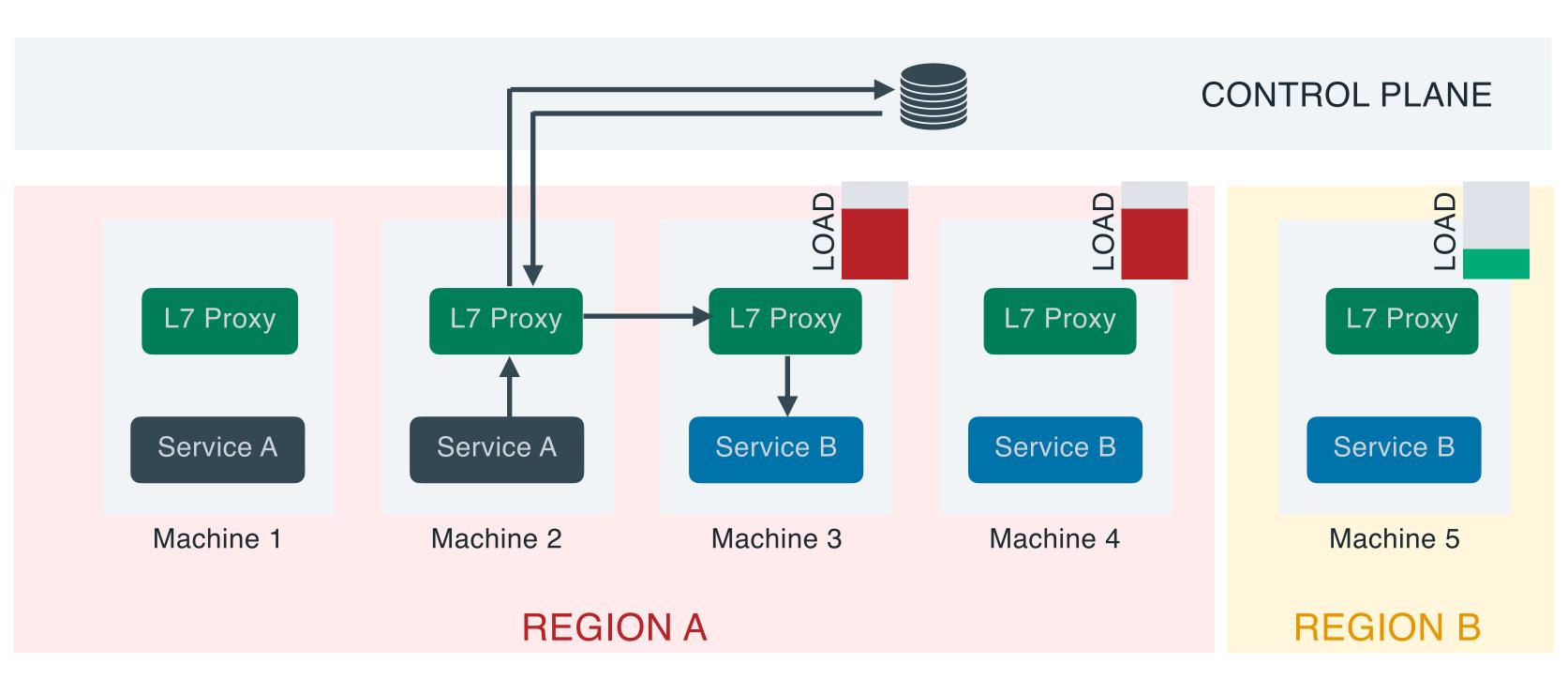
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 - P90 cross-region latency: 106ms





Service Mesh Challenges

- [SCALABILITY] How can we scale service discovery to O(10⁶) clients and proxies?
- [HW COST] How to minimize HW cost?
- [RPC LATENCY & LB] How to simultaneously minimize RPC latency and load balance across geo-distributed hosts?
 - Sidecars add extra latency
 - O(10-10⁴) hosts per service
 - P90 cross-region latency: 106ms
- [SHARDED SERVICES] Support for shared services NOT COVERED



03 ServiceRouter

KEY DESIGN CONCEPTS

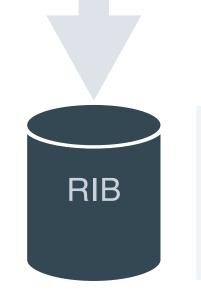
RIB

Routing Information Base

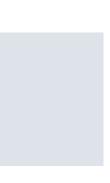
Decentralize the unscalable part of the control plane in order to scale out.

 Independent controllers execute different functions such as registering services and generating a per-service cross-region routing table.

CONTROLLERS



- Service Discovery Info
- Per-service routing config
- Cross-region service routing info

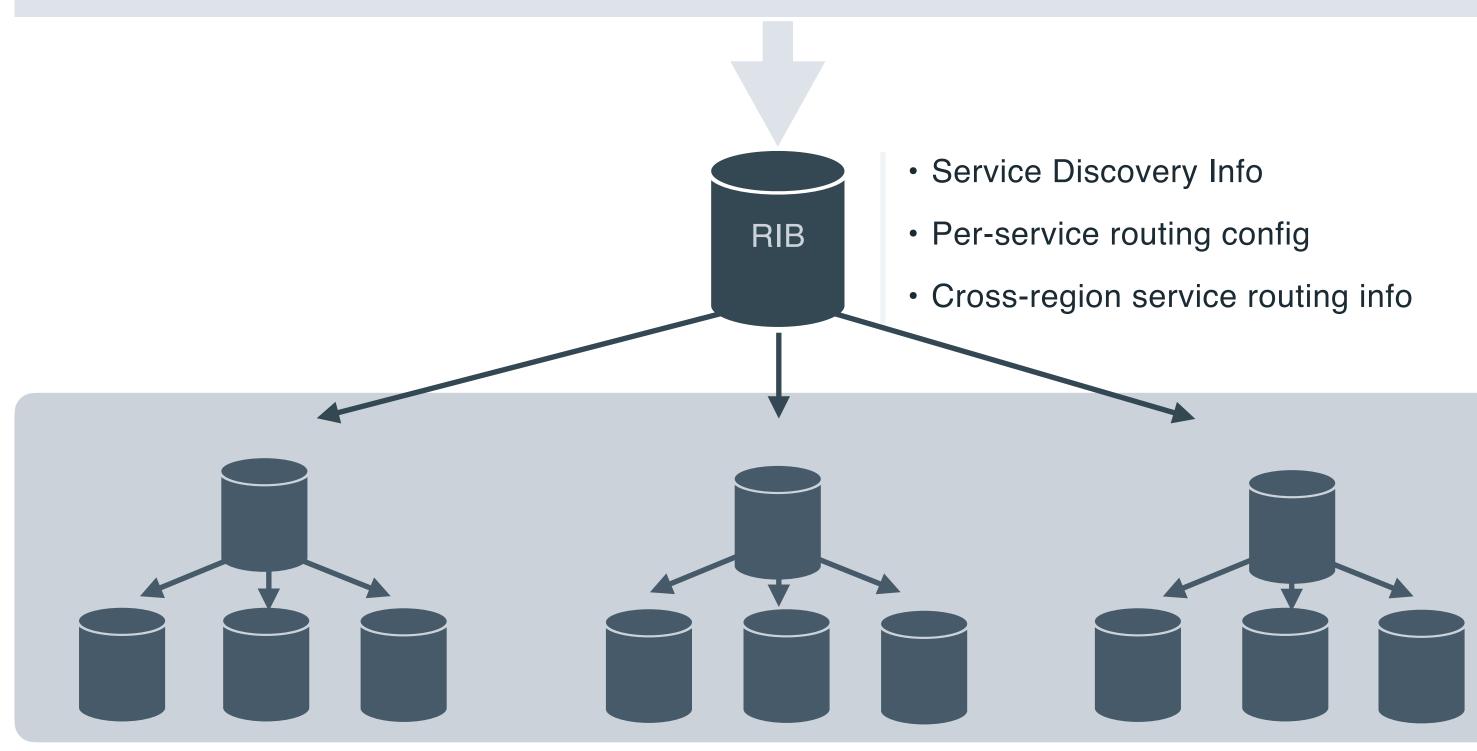


RIB

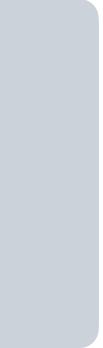
Routing Information Base

Decentralize the unscalable part of the control plane in order to scale out.

- Independent controllers execute different functions such as registering services and generating a per-service cross-region routing table.
- The data distribution layer massively replicates the RIB so that there are sufficient RIB replicas to handle read traffic from millions of proxies.



CONTROLLERS



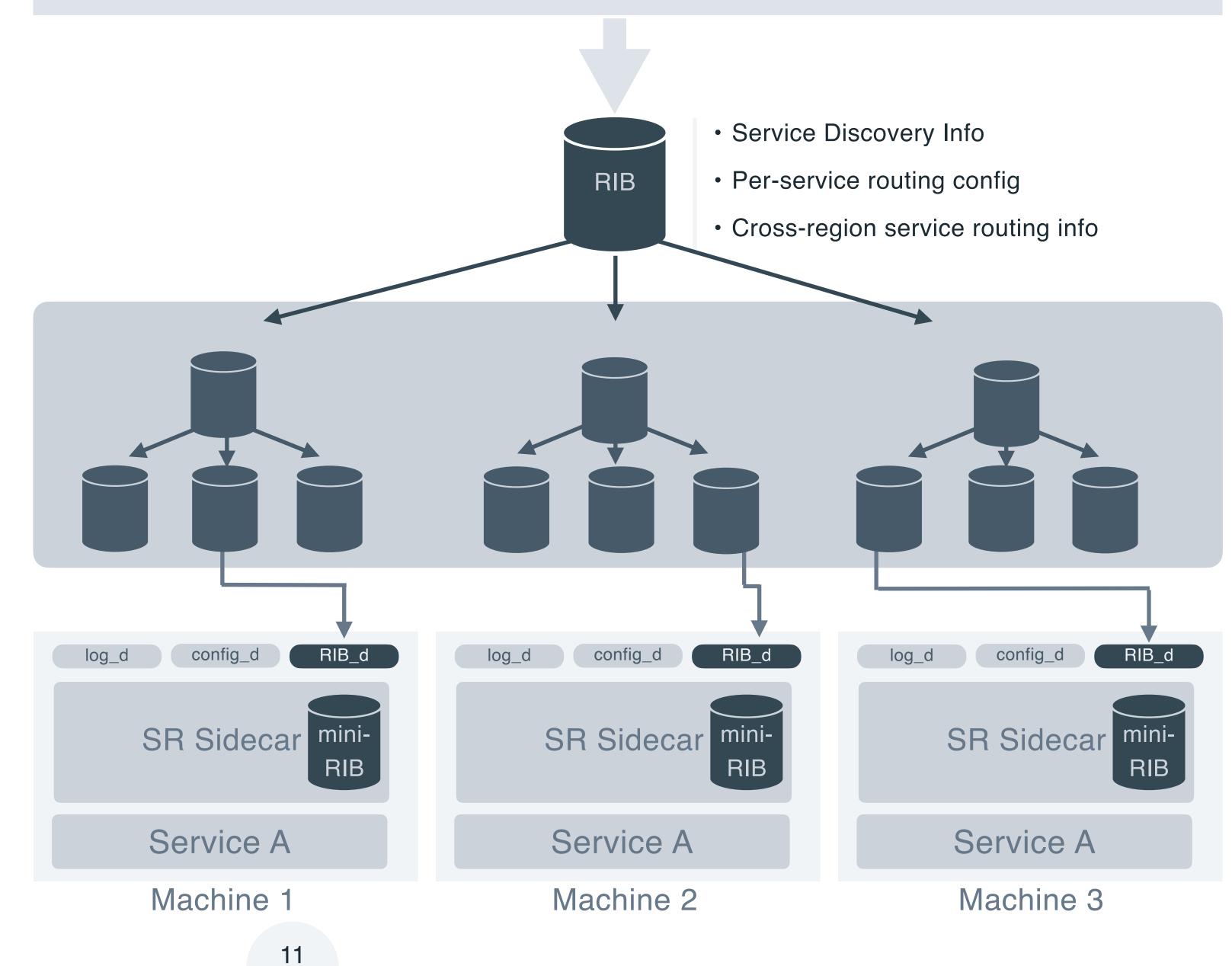
Data Distribution Layer

RIB

Routing Information Base

Decentralize the unscalable part of the control plane in order to scale out.

- Independent controllers execute different functions such as registering services and generating a per-service cross-region routing table.
- The data distribution layer massively replicates the RIB so that there are sufficient RIB replicas to handle read traffic from millions of proxies.
- Each proxy self-configures and selfmanages without the control plane's direct involvement.

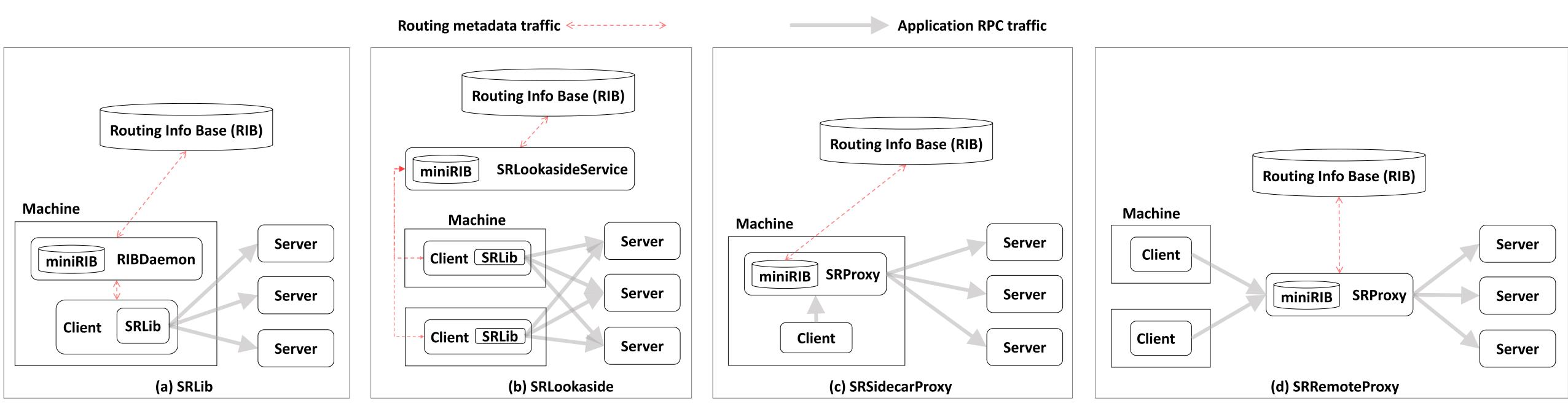


CONTROLLERS

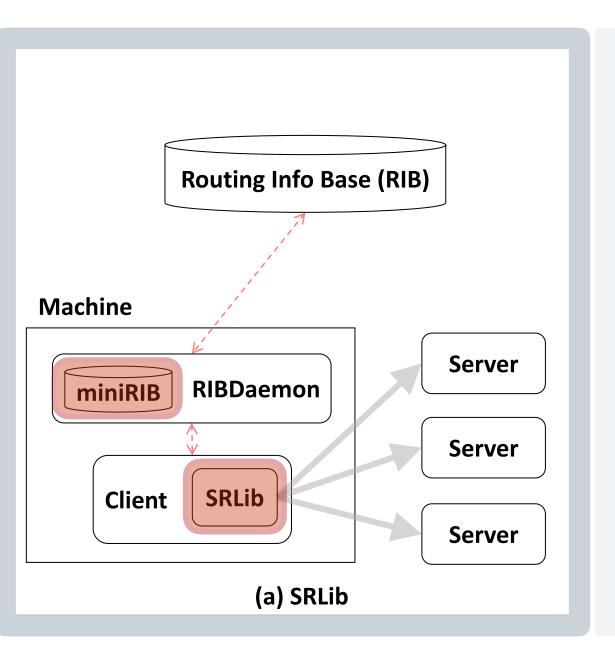


Versatility

Controllers are agnostic to the L7 architecture.



RPC traffic routed 99% through SRLib.



Routing metadata traffic <----->

SRLib

Provide the service-mesh functions out of a library that is directly linked into the RPC client's executable

• Eliminates side car latency overhead

Run a separate RIBDaemon on the client machine to cache miniRIB.

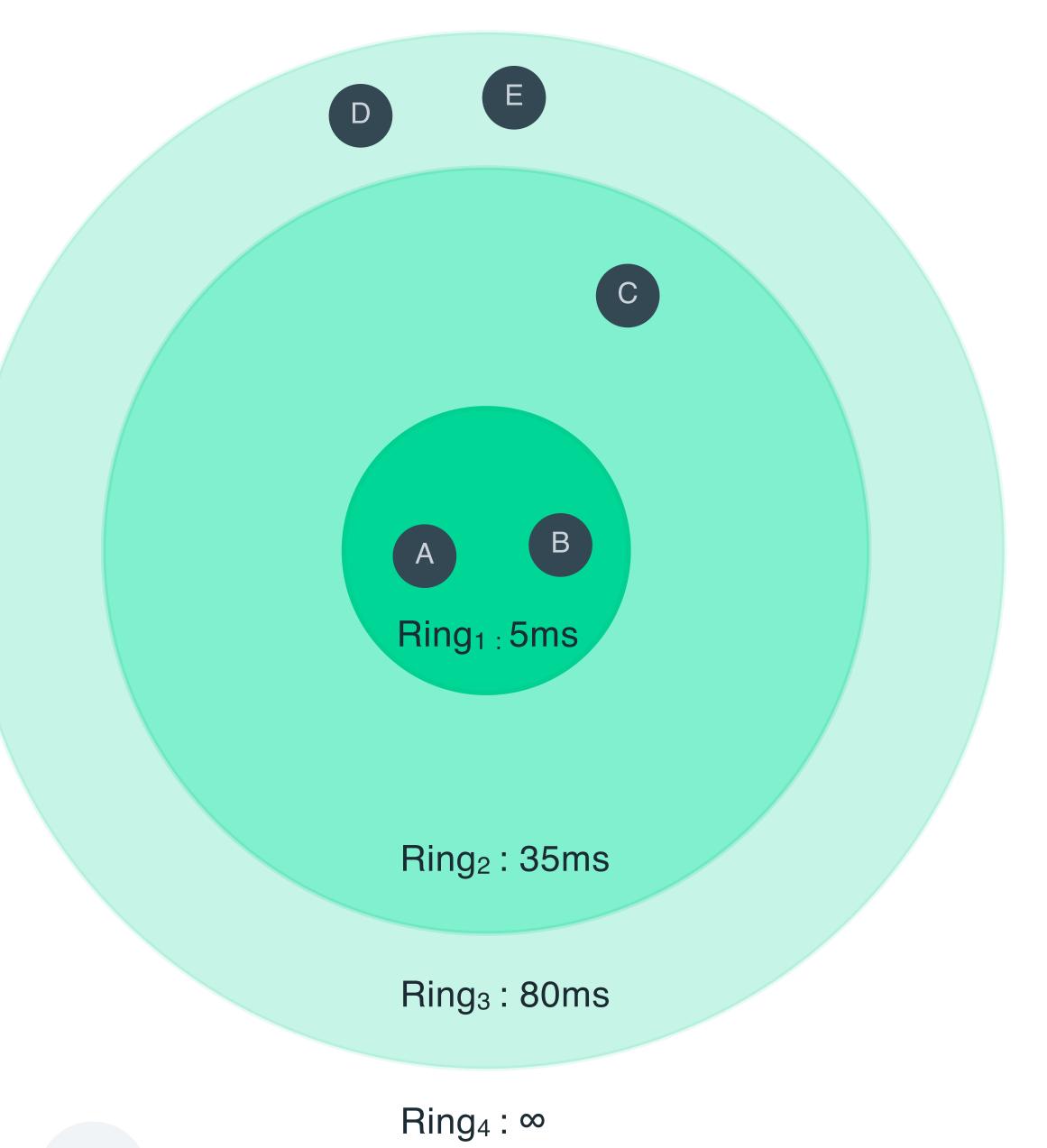


 $Ring_1: 5ms \ I \ Ring_2: 35ms \ I \ Ring_3: 80ms \ I \ Ring_4: \infty$

LATENCY RINGS AND CROSS-REGION ROUTING

SR strives to simultaneously minimize RPC latency and balance load across global regions.

- SR introduces the concept of latency rings to minimize latency.
- SR collects per-service global traffic and load information, computes a per-service cross-region routing table, and disseminate it to L7 routers to guide their local routing decisions.



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Ring₁ : 5ms | Ring₂ : 35ms | Ring₃ : 80ms | Ring₄ : ∞

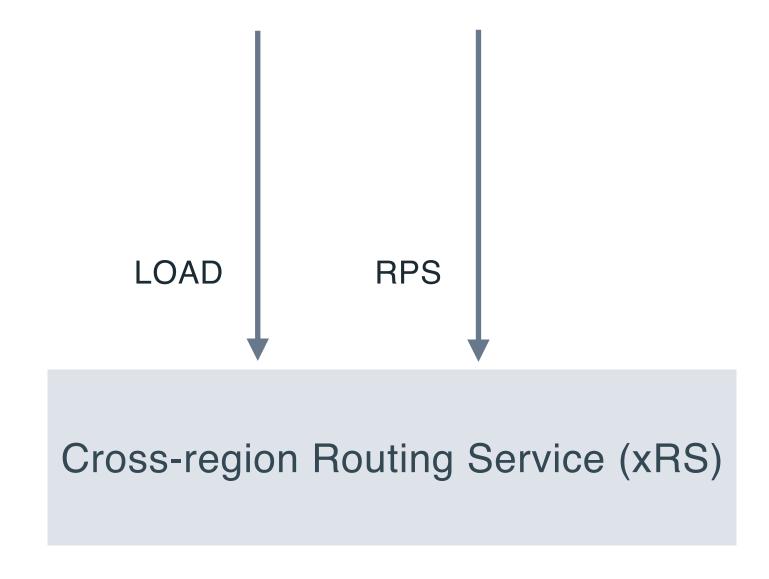


Ring₄ : ∞

LATENCY RINGS AND CROSS-REGION ROUTING

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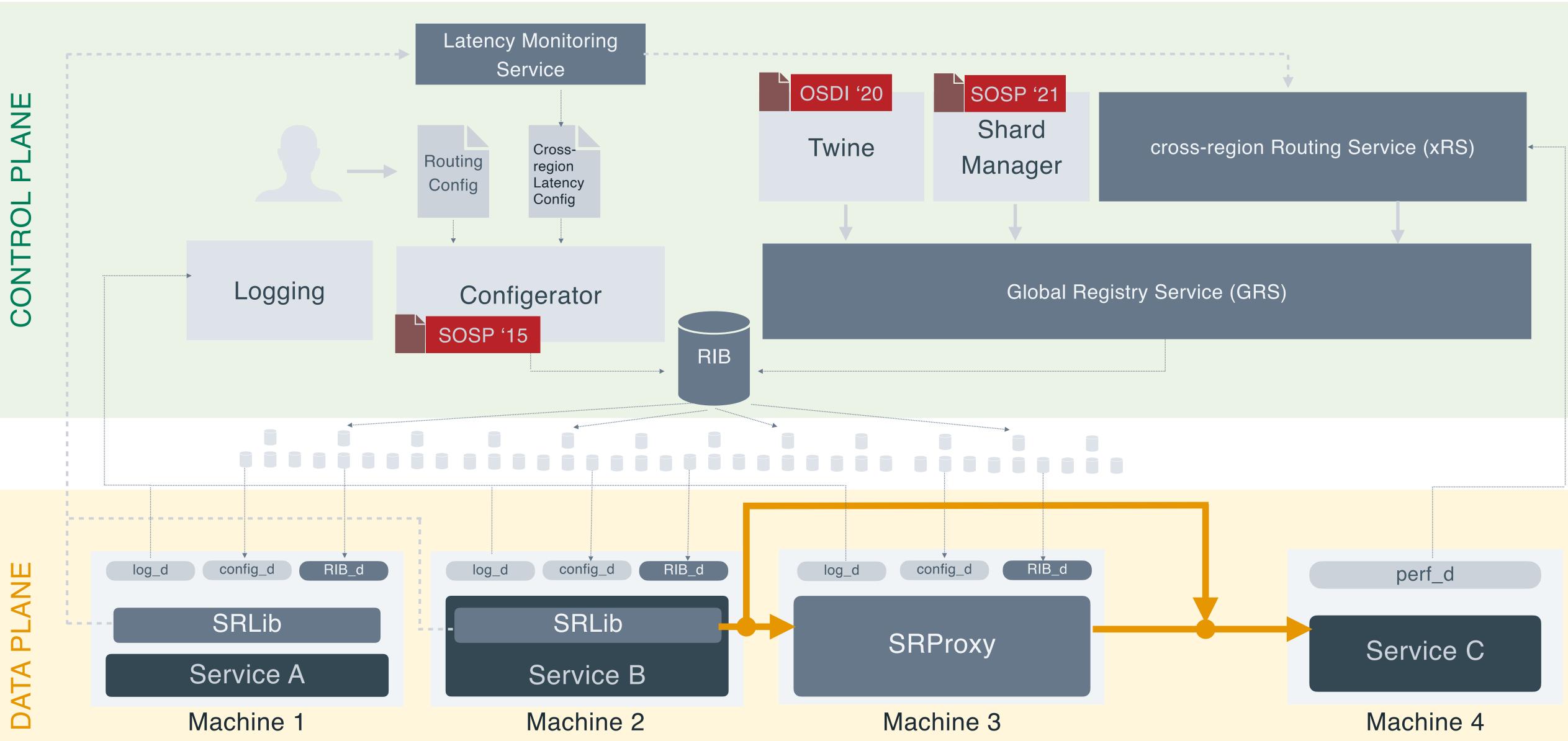
Ring₁ : 5ms : 55% | Ring₂ : 35ms : 65% | Ring₃ : 80ms : 80% | Ring₄ : ∞ : ∞

Load threshold for Ring₁

04 ServiceRouter

OVERALL ARCHITECTURE

ServiceRouter Architecture 04



Ω

05 ServiceRouter

EVALUATION

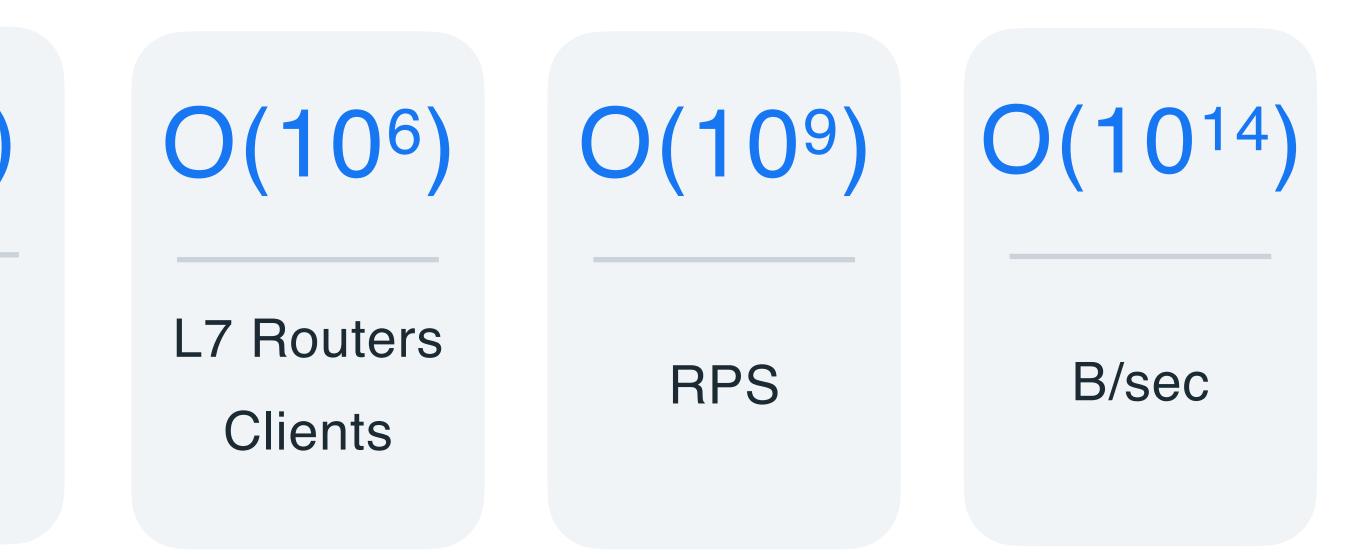
Scalability

Overall scale

- Regions
- Routers/Clients/Servers
- Throughput

O(10)

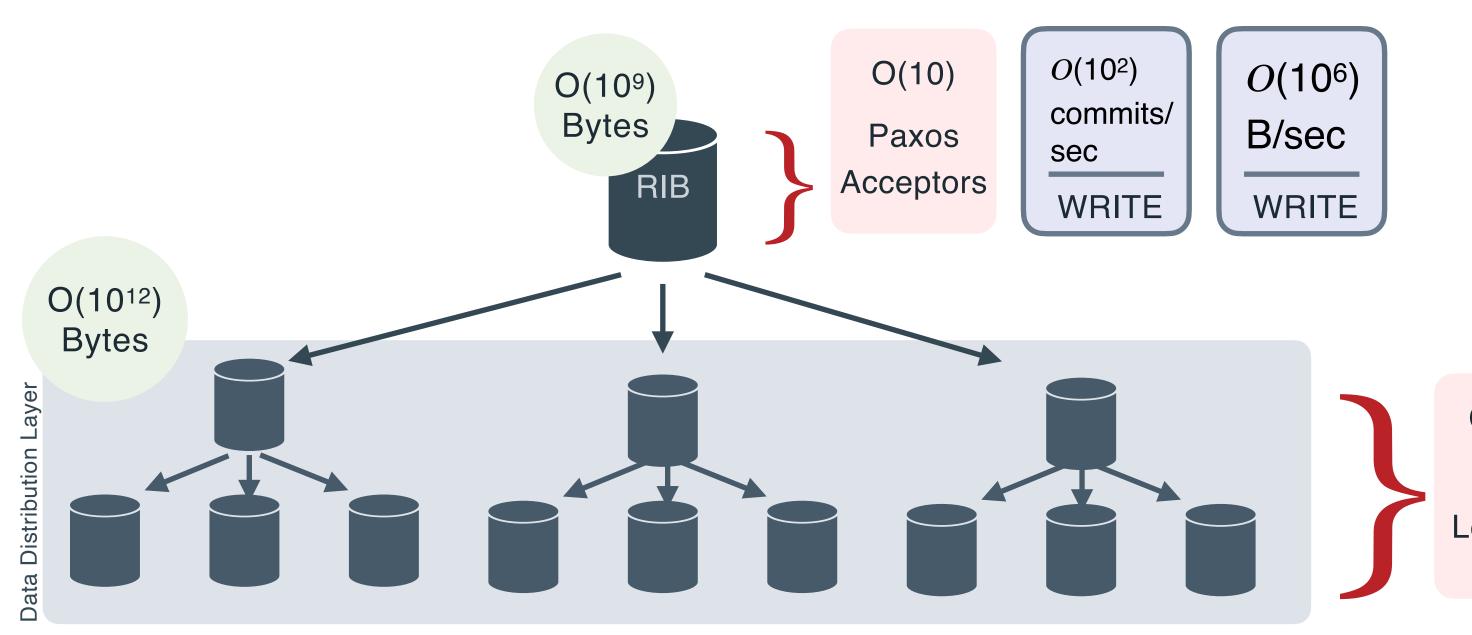
Regions



Scalability

RIB - Routing Information Base

- RIB Replicas
- RIB Write bandwidth
- RIB Write throughput



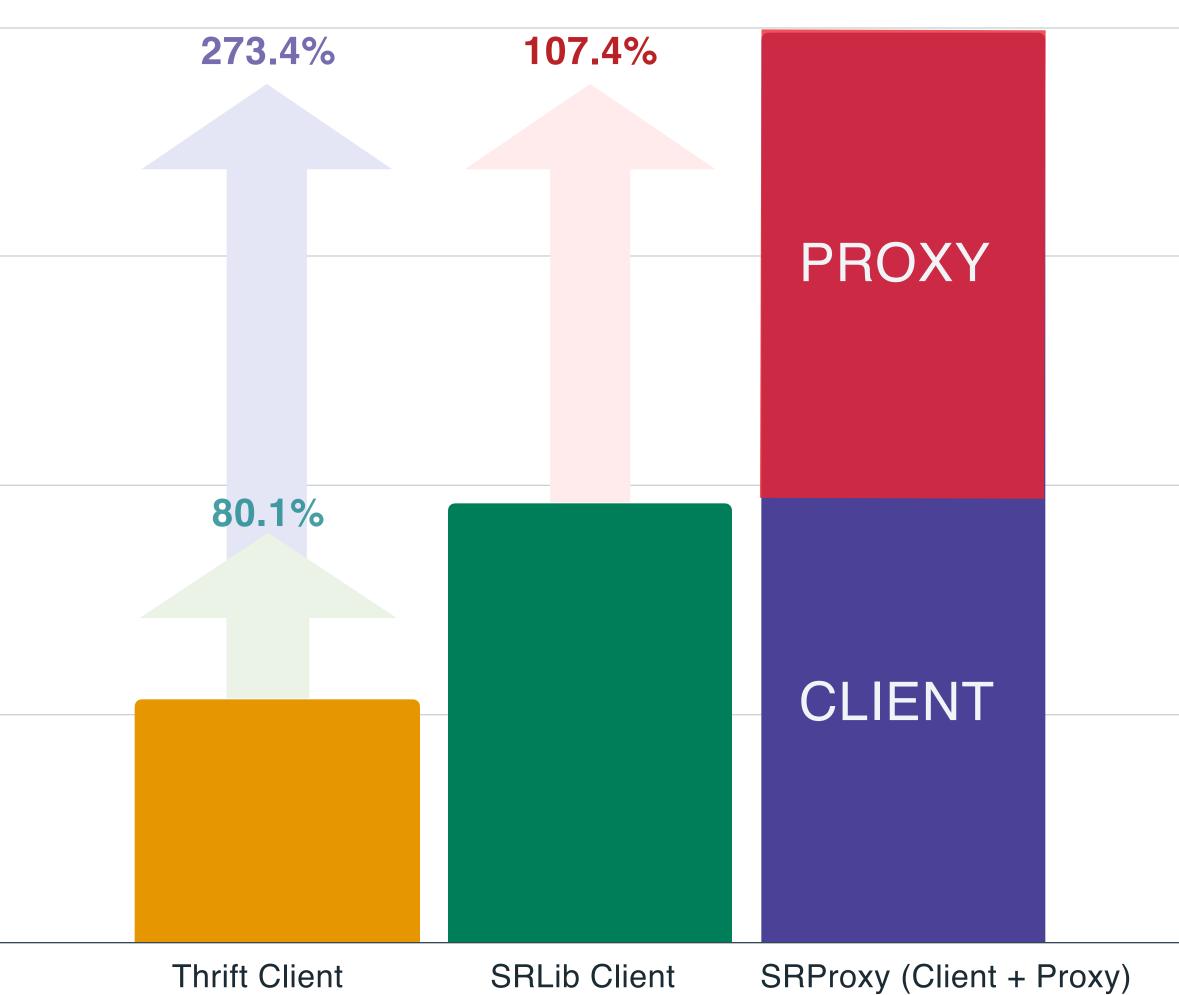


Cost

METHODOLOGY

•	Metrics: P50 avg request latency; CPU Instructions per request	450000 —	
•	Designs		
	- Baseline: Thrift RPC		
	- SRLib		
	- Remote SRProxy	300000	
•	Simulated Payload:		
	 Production avg request and avg response size 		
	- O(10 ³) B	150000 —	
•	100K requests		
•	3 trials per design		

CPU Instructions/Request

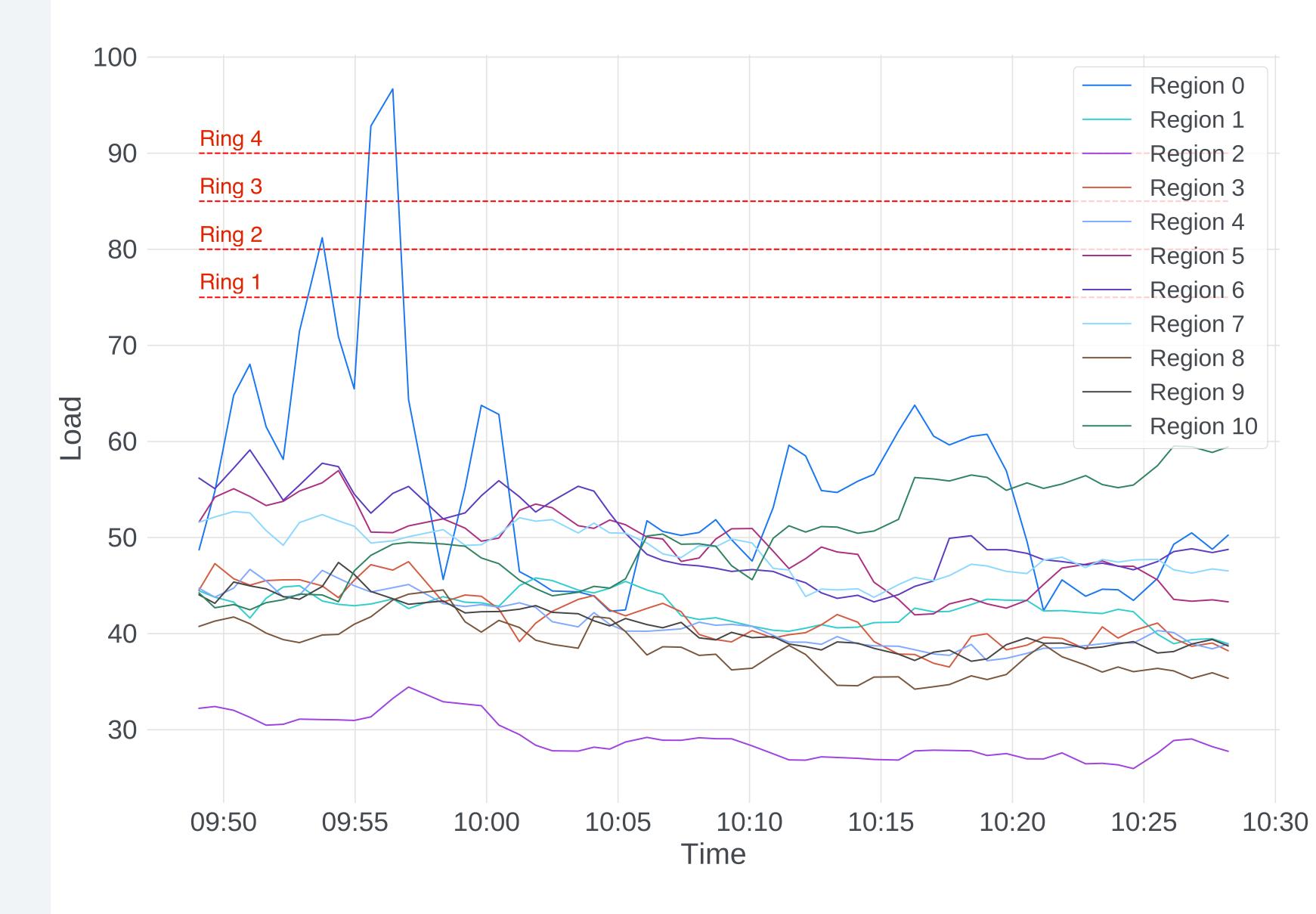


0

600000

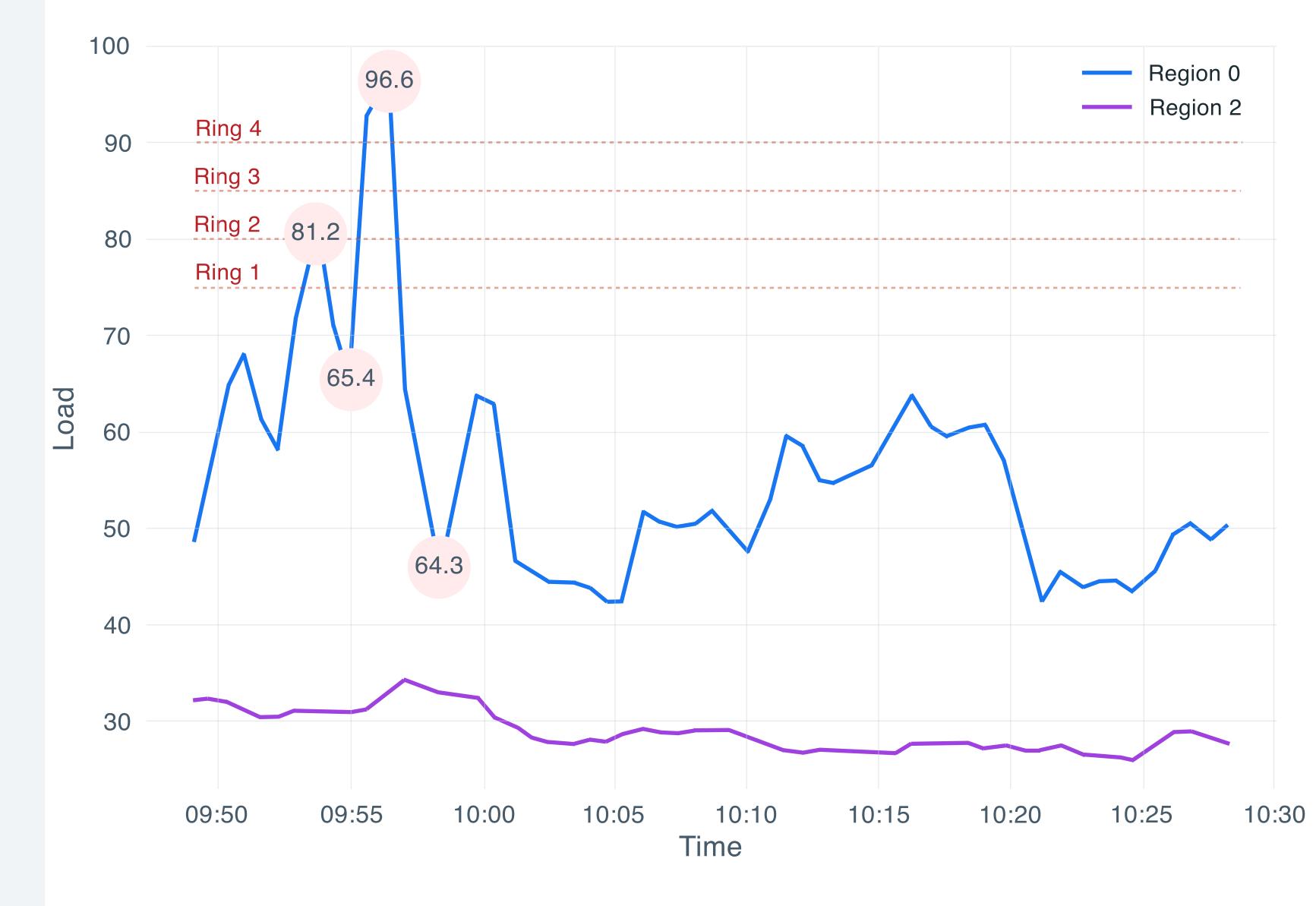
Cross-Region Load Shift

Real-world Example



Cross-Region Load Shift

- Real-world Example
- 9:53 —> Region 0 Load = 81.2%
 - xRS Traffic Shifts
 - R0 -5.35%
 - R0 to R2 +5.35%
- 9:54 —> Region 0 Load = 65.47%
- 9:56 —> Region 0 Load = 96.69%
 - xRS Traffic Shifts
 - R0 -25%
 - R0 to R2 +25%
- 9:57 —> Region 0 Load = 64.34%



ServiceRouter

HYPERSCALE AND MINIMAL COST SERVICE MESH AT META

Summary 06



Imperial College London



ServiceRouter's massive RIB replication allows decentralizing L7 router management and to scale to millions of routers and proxies.

ServiceRouter routes 99% of the traffic with an optimized embedded library approach with astounding HW savings.

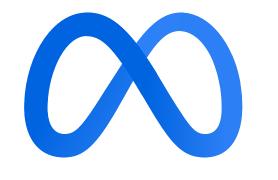
ServiceRouter's source-based locality rings and xRS strike a balance between latency wins and load balancing.

Built-in **support for sharded services** which account for 68% of our RPCs [not covered in this talk].











Soteris Demetriou I <u>s.demetriou@imperial.ac.uk</u>

